

## NCU Microeconomics: classnote 4

*Mechanism Design II* : Groves Mechanism and AGV Mechanism

- $I + 1$  players.

- Types are drawn from independent distribution  $P_i$  on  $[\underline{\theta}_i, \bar{\theta}_i]$  with strictly positive and differentiable densities  $p_i$ . The distributions are common knowledge.
- (private values) Agent  $i$ 's preferences:  $u_i(x, t_i, \theta_i)$
- Preferences are quasi-linear: For any  $i \in \{1, \dots, I\}$ ,

$$u_i(x, t_i, \theta_i) = V_i(x, \theta_i) + t_i,$$

and either

$$u_0(x, t, \theta) = V_0(x, \theta) - \sum_{i=1}^I t_i$$

(self-interested principal) or

$$u_0(x, t, \theta) = \sum_{i=0}^I V_i(x, \theta)$$

(benevolent principal), where  $V_0(x, \theta) = B_0(x, \theta) - C_0(x)$ ,  $C_0(x)$  is the principal's monetary cost from decision  $x$  and  $B_0(x, \theta)$  is nonmonetary benefit.

- An allocation  $y(\cdot)$  is (*ex post*) *efficient* if  $x(\theta) \in K$  for each  $\theta$  and

$$x(\theta) \text{ maximizes } \sum_{i=0}^I V_i(x, \theta) \text{ over } K, \text{ for all } \theta$$

- Budget Balance:

$$\sum_{i=1}^I t_i(\theta) \leq -C_0(x(\theta)) \text{ for all } \theta$$

- Dominant Strategy vs. Bayesian Mechanisms: Choose agent  $i$ 's transfer so that agent  $i$ 's payoff is the same as the total surplus of all parties up to a constant.

- Dominant-strategy mechanism: Each agent's optimal announcement is independent of the announcements of the other agents, i.e., IC is for each agent  $i = 1, \dots, I$  and for each  $\theta_i, \hat{\theta}_i$  and  $\theta_{-i}$ ,

$$u_i(y(\theta_i, \theta_{-i}), \theta_i) \geq u_i(y(\hat{\theta}_i, \theta_{-i}), \theta_i).$$

- Bayesian mechanism: IC is

$$E_{\theta_{-i}} u_i(y(\theta_i, \theta_{-i}), \theta_i) \geq E_{\theta_{-i}} u_i(y(\hat{\theta}_i, \theta_{-i}), \theta_i)$$

- Dominant strategy mechanism is not sensitive to beliefs that players have about each other and it does not require players to compute Bayesian equilibrium strategies. However, focusing on dominant-strategy mechanism restricts the set of mechanisms considerably. Thus, implementation in dominant strategies is a nice property to have if feasible, but it is not clear how much utility loss a principal should be willing to tolerate in order to have dominant strategies for the agents.
- Mookherjee and Reichelstein (1989) identify a class of models in which dominant strategy implementation involves no welfare loss relative to Bayesian implementation.

- Efficiency Theorems

- The Groves Mechanism (1973) Dominant Strategy implementation: Choose agent  $i$ 's transfer so that agent  $i$ 's payoff is the same as the total surplus of all parties up to a constant.

- \* Let  $x^*(\theta)$  maximizing  $\sum_{i=0}^I V_i(x, \theta_i)$  denote an efficient solution for the type profile  $\theta$ .

- \* Define  $t_i(\hat{\theta}) \equiv \sum_{\substack{j \in \{0, \dots, I\} \\ j \neq i}} V_j(x^*(\hat{\theta}_i, \hat{\theta}_{-i}), \hat{\theta}_j) + \tau_i(\hat{\theta}_{-i})$ , where  $\tau_i(\hat{\theta}_{-i})$  is an arbitrary function of  $\hat{\theta}_{-i}$ .

- \* Show that it is optimal for agent  $i$  to announce his true type ( $\hat{\theta}_i = \theta$ ) regardless of the other agents' announcements. This implies that  $(x^*(\hat{\theta}), t(\hat{\theta}))$  is a dominant strategy mechanism which yields efficient allocation. The proof is simple: Suppose that agent  $i$  strictly prefers announcing  $\hat{\theta}_i$  to announcing  $\theta_i$  for some types  $\hat{\theta}_{-i}$  of

the other agents. Then

$$\begin{aligned} & V_i \left( x^* \left( \hat{\theta}_i, \hat{\theta}_{-i} \right), \theta_i \right) + \sum_{j \neq i} V_j \left( x^* \left( \hat{\theta}_i, \hat{\theta}_{-i} \right), \hat{\theta}_j \right) \\ & > V_i \left( x^* \left( \theta_i, \hat{\theta}_{-i} \right), \theta_i \right) + \sum_{j \neq i} V_j \left( x^* \left( \theta_i, \hat{\theta}_{-i} \right), \hat{\theta}_j \right). \end{aligned}$$

But this contradicts the fact that  $x^* \left( \theta_i, \hat{\theta}_{-i} \right)$  is efficient for type profile  $\left( \theta_i, \hat{\theta}_{-i} \right)$ .

\* Budget balance may not be satisfied.

\* Example: Should a bridge be built?  $u_i = \theta_i x + t_i$ , where  $x$  is equal to 0 or 1 and  $\theta_i$  is agent  $i$ 's valuation or willingness to pay for the public good. With  $c > 0$  denoting the cost of supplying the public good, the efficient rule is

$$x^* (\theta) = \begin{cases} 1 & \text{if } \sum_{i=1}^I \theta_i \geq c \\ 0 & \text{otherwise} \end{cases}$$

One Groves mechanism for this example takes the following form:

$$t_i (\hat{\theta}) = \begin{cases} \sum_{j \neq i} \hat{\theta}_j - c & \text{if } \sum_{j=1}^I \hat{\theta}_j \geq c \\ 0 & \text{otherwise} \end{cases}$$

and

$$x^* (\hat{\theta}) = \begin{cases} 1 & \text{if } \sum_{i=1}^I \hat{\theta}_i \geq c \\ 0 & \text{otherwise} \end{cases}$$

– The AGV Mechanism: Instead of being paid the surpluses of the other agents on the basis of their reports, each agent is paid the expected value of the other agents' surpluses conditional on his own report.

$$t_i (\hat{\theta}) \equiv E_{\theta_{-i}} \left( \sum_{\substack{j \in \{0, \dots, I\} \\ j \neq i}} V_j \left( x^* \left( \hat{\theta}_i, \theta_{-i} \right), \theta_j \right) \right) + \tau_i (\hat{\theta}_{-i})$$

The function  $\tau_i (\cdot)$  will be determined later to ensure BB.

\*  $(x^*, t)$  is (Bayesian) incentive compatible:  $\hat{\theta}_i = \theta_i$  must maximize

$$E_{\theta_{-i}} \left( V_i \left( x^* \left( \hat{\theta}_i, \theta_{-i} \right), \theta_i \right) + \sum_{j \neq i} V_j \left( x^* \left( \hat{\theta}_i, \theta_{-i} \right), \theta_j \right) \right).$$

But  $\hat{\theta}_i = \theta_i$  maximizes the term inside the expectation operator for all  $\theta_{-i}$  and, therefore, maximizes the expectation.

\* Suppose cost  $C_0(x) = 0$ . BB requires  $\sum_{i=1}^I t_i(\hat{\theta}) = 0$ . Let

$$\mathcal{E}_i(\hat{\theta}_i) \equiv E_{\theta_{-i}} \left( \sum_{j \neq i} V_j(x^*(\hat{\theta}_i, \theta_{-i}), \theta_j) \right)$$

denote the “expected externality” for agent  $i$  when he announces  $\hat{\theta}_i$ . Since  $\mathcal{E}_i(\hat{\theta}_i)$  is the first part of the transfer to agent  $i$  and  $\tau_i(\cdot)$  is supposed not to depend on  $\hat{\theta}_i$ ,  $\mathcal{E}_i(\hat{\theta}_i)$  must be paid by other agents, i.e.,

$$\begin{aligned} \tau_i(\hat{\theta}_{-i}) &= - \sum_{j \neq i} \frac{\mathcal{E}_j(\hat{\theta}_j)}{I-1} \\ &= - \frac{1}{I-1} \sum_{j \neq i} E_{\theta_{-j}} \left( \sum_{k \neq j} V_k(x^*(\hat{\theta}_j, \theta_{-j}), \theta_k) \right). \end{aligned}$$

\* For  $C_0(x) > 0$ . BB requires  $\sum_{i=1}^I t_i(\hat{\theta}) \leq -C_0(x(\hat{\theta}))$ . Consider the “fictional problem” where the agents’ utility functions are

$$\tilde{V}_i(x, \theta_i) \equiv V_i(x, \theta_i) - \frac{C_0(x)}{I}$$

and the principal’s cost is  $\tilde{C}_0(x) \equiv 0$ . We then compute the transfers  $\tilde{t}_i(\cdot)$  for this fictional problem, and set

$$t_i(\cdot) = \tilde{t}_i(\cdot) - C_0(x^*(\hat{\theta}))/I.$$

\* Ex post IR may not be satisfied. Need agents to sign a contract before they learn their types privately.