

DESIGNING RANDOM ALLOCATION MECHANISMS:
DISCUSSION

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KEY ASSUMPTIONS ABOUT THE ENVIRONMENT

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- ▶ Resale can be fully controlled
- ▶ Many agents
- ▶ Agents should be treated symmetrically – hence reliance on randomization

GOOD ORDINAL MECHANISMS

- ▶ Strategy-proof
- ▶ Ordinally efficient
- ▶ Symmetric (equal treatment of equals)

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- ▶ Probabilistic Serial (Bogomolnaia and Moulin 2001)

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Che and Kojima 2010:

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Core from Random Endowments (Top Trading Cycles) coincide with Random Priority already in small markets (Abdulkadiroglu and Sonmez 1999, Pathak and Sethuraman 2011)

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Theorem (Liu and Pycia 2011). In the continuum economy, every mechanism that is strategy-proof, ordinally efficient, and symmetric coincides with Random Priority/Probabilistic Serial for almost every preference distribution.

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If we restrict attention to "continuous" mechanisms then:

- ▶ In the continuum economy, Random Priority/Probabilistic Serial is the only strategy-proof, ordinally efficient, and symmetric mechanism.
- ▶ If a sequence of symmetric and strategy-proof mechanisms is asymptotically OE as we replicate the economy, then the sequence converges to Random Priority/Probabilistic Serial.

INTENSIVE VS EXTENSIVE MARGINS

Manea 2009: Random Priority does not become ordinally efficient in the limit as the number of object types grows.

ORDINAL AND CARDINAL MECHANISMS

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- ▶ Cardinal mechanisms offer the premise of eliciting preference intensity, and hence implement Pareto-efficient outcomes

TRADE-OFFS IN DESIGNING CARDINAL MECHANISMS

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- ▶ Boston mechanism: good efficiency properties in equilibrium (Abdulkadiroglu, Che, and Yasuda 2011) but fails strategy-proofness even in the limit (Abdulkadiroglu and Sonmez 2003, Kojima and Pathak 2009) and has bad redistributive properties (Pathak and Sonmez 2008)

HOW TO DECOMPOSE A RANDOM ALLOCATION?

Budish, Che, Kojima, Milgrom 2010 show how to decompose allocations while preserving conjunctions of elementary constraints of the form

$$\underline{c} \leq \sum_{(i,h) \in C} P(i,h) \leq \bar{c}$$

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