

Let (W, S) Coxeter group. $W = \langle S \mid \underbrace{s_i^2 = \dots = s_{|S|}^2 = \text{Id}}_{\text{(quadratic rels)}}, \underbrace{\{sts\} \dots = \{tst\} \dots}_{\substack{\text{mst} \\ \text{(braid rels)}}} \rangle$

Given $w \in W$,
 $l(w) :=$ minimum word length in S 's
 When w achieved $l(w)$, called w reduced expr.

[4] (Matsumoto) All reduced words are related by braid rels.

[5] (Exchange Condition)

Given $w = s_1 s_2 \dots s_k$ reduced expr. Given $s \in S$.
 If $l(ws) < l(w)$ then $\exists i \leq k$ s.t. $ws = s_1 s_2 \dots \hat{s}_i \dots s_k$.

[6] (Useful) In above case, $w = wss = s_1 \dots \hat{s}_i \dots s_k s$ is reduced & ending at s
i.e. $\forall s \in S$, either $l(ws) = k+1$ or $\exists i$ s.t. $w = s_1 \dots \hat{s}_i \dots s_k s$

[7] (Deletion Condition)

Given $x = s_1 \dots s_k$ any express
 If $l(x) < k$, then $\exists i < j < k$ s.t. $x = s_1 \dots \hat{s}_i \dots \hat{s}_j \dots s_k$.

[8] (Remark) In view of exchange condition, not every s_j can be eliminated by $(\) \circ s$

However, can be achieved by $(\) \circ t \in T$, $T = \{t \in W \mid \exists \text{ some } s \in S\}$
 Namely, let $t = s_k \dots s_{j+1} s_j s_{j+1} \dots s_k$. \forall *T are all reflections!*

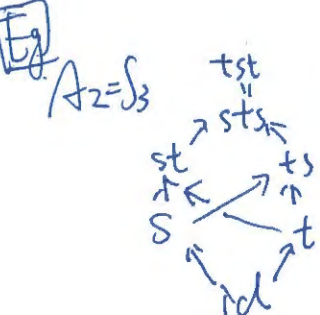
Then $wt = s_1 \dots \hat{s}_j \dots s_k$ eliminates s_j for arbitrary j

[9] (Bruhat order)

We write $x \rightarrow y$ if $\begin{cases} l(x) < l(y) \\ x = yt \text{ (equiv, } xt = y) \text{ for some } t \in T. \end{cases}$

[10] (Bruhat order)

$w \leq z$ if $\exists w \rightarrow a \rightarrow b \rightarrow \dots \rightarrow y \rightarrow z$
 $x \leq y$ means some substring of some reduced exprs of y is a reduced exprs of x .



[11] (Remark) It does not matter to use left or right T -multiplication in the def. of Bruhat order
 $\therefore x = yt = \underbrace{(yt^{-1})}_{\substack{\text{mst} \\ \text{rels}}} y$ *(Bruhat order is intrinsic)*
(but for Coxeter cplx always on the right!)

Given $W = (W, S) \rightsquigarrow \mathbb{Z}[W]$ its group alg $\xrightarrow{\text{deform}}$ $H = H(W)$
 $\xleftarrow{V=1}$ Hecke alg over $\mathbb{Z}[v, v^{-1}]$

Hecke alg

H is the unital ($1 = \delta_{id}$)

assoc. alg. $\otimes \mathbb{Z}[v, v^{-1}]$ gen. by $\{\delta_s / s \in S\}$ w/ relations

(quadratic rels) $\delta_s^2 = (v^{-1} - v)\delta_s + 1$ may think then $v^{-1} - v \in i\mathbb{R}$.
 $v \in S^1 \subset \mathbb{C}$ only when $v=1$ gives $0 \in \mathbb{R}$.

(braid rels) $\delta_{s_1} \delta_{s_2} \delta_{s_1} \dots = \delta_{s_2} \delta_{s_1} \delta_{s_2} \dots$ by Matsumoto, well-defined

Standard Basis

$\forall x \in W, x = s_1 s_2 \dots s_m$ reduced \otimes define $\delta_x := \delta_{s_1} \delta_{s_2} \dots \delta_{s_m} \in H$

Standard basis of Hecke alg $H: \{\delta_x / x \in W\}$

Product Rule

$$\delta_x \delta_s = \begin{cases} \delta_{xs} & x < xs \\ (v^{-1} - v)\delta_x + \delta_{xs} & x > xs \end{cases} \quad \delta_s \delta_x = \begin{cases} \delta_{sx} & x < sx \\ (v^{-1} - v)\delta_x + \delta_{sx} & x > sx \end{cases}$$

$x = s_1 \dots s_n$ reduced expr. if $x < xs \Rightarrow xs = s_1 \dots s_n s$ is reduced $\Rightarrow \delta_x \delta_s = \delta_{s_1} \dots \delta_{s_n} \delta_s$
 if $x > xs$, by useful cor rewrite $x = s_1 \dots s_{n-1} s$ & $xs = s_1 \dots s_{n-1}$ reduced. $= \delta_{xs}$
 $\Rightarrow \delta_x \delta_s = \delta_{s_1} \dots \delta_{s_{n-1}} \delta_s \delta_s = \delta_{s_1} \dots \delta_{s_{n-1}} [(v^{-1} - v)\delta_s + 1] = (v^{-1} - v)\delta_x + \delta_{xs}$

Inversion Rule

$\delta_s^2 = (v^{-1} - v)\delta_s + 1 \Rightarrow \delta_s = (v^{-1} - v) + \delta_s^{-1} \Rightarrow \delta_s^{-1} = \delta_s + (v - v^{-1})$ algebra when it comes to inverse, the product is opposite to group product

Let $w = s_1 \dots s_m$ reduced. $\Rightarrow w^{-1} = s_m \dots s_1 \Rightarrow \delta_{w^{-1}} = \delta_{s_m} \dots \delta_{s_1} \Rightarrow (\delta_{w^{-1}})^{-1} = \delta_{s_1}^{-1} \dots \delta_{s_m}^{-1}$
 $\Rightarrow (\delta_{w^{-1}})^{-1} = [\delta_{s_1} + (v - v^{-1})] \dots [\delta_{s_m} + (v - v^{-1})] = \sum_{e \in \{0,1\}^m} A_{e(v^{-1}v)} \cdot \delta_{s_1}^{e_1} \dots \delta_{s_m}^{e_m}$

Want to show $\delta_{s_1}^{e_1} \dots \delta_{s_m}^{e_m} \in \sum_{x \in W} \mathbb{Z}[v, v^{-1}] \delta_x$

Write $\delta_{s_1}^{e_1} \dots \delta_{s_m}^{e_m} = \delta_{r_1} \delta_{r_2} \dots \delta_{r_k}$ by picking those $e_j = 1$.
 if not reduced, let $r_1 r_2 \dots r_j$ reduced $\Rightarrow (r_1 r_2 \dots r_j) r_{j+1}$ not reduced.
 By Deletion Corollary $\exists 1 \leq a < b \leq j+1$ st. $r_1 r_2 \dots r_{j+1} = r_1 \dots r_a \dots r_b \dots r_{j+1}$ reduced.

\Rightarrow By Product Rule $(\delta_{r_1} \dots \delta_{r_{j+1}}) \delta_{r_{j+2}} \dots \delta_{r_k} = [(v^{-1} - v)\delta_{r_1 \dots r_j} + \delta_{r_1 \dots r_a \dots r_b \dots r_{j+1}}] \delta_{r_{j+2}} \dots \delta_{r_k}$

$$= (\prod v) \delta_{i_1} \dots \delta_{j_{i_1}} \dots \delta_k + \delta_{i_1} \dots \delta_{j_{i_2}} \dots \delta_k \dots \delta_k$$

again, choose $\max \{i \text{ s.t. } \dots\}$ $\max \{i \text{ s.t. } \dots\}$

repeat the procedure we finally get reduced expr. for each terms,

therefore $\delta_{s_1}^{e_1} \dots \delta_{s_m}^{e_m} \in \sum_{x \subseteq W} \mathbb{Z}[v, v^{-1}] \circ \delta_x$.

So we prove

Prop The standard basis δ_W is invertible and.

$$\delta_W^{-1} = \delta_W + \sum_{x \subset W} a_x \cdot \delta_x, \quad a_x \in \mathbb{Z}[v, v^{-1}]$$

Kazhdan-Lusztig Basis

Observe that $\delta_s^{-1} = \delta_s + (v - v^{-1}) \Rightarrow \delta_s^{-1} + v^{-1} = \delta_s + v$ } self dual. } $b_s := \delta_s + v$ is

Define bar involution $H \xrightarrow{\bar{\cdot}} H$ by $\bar{v} = v^{-1}, \bar{\delta}_s := \delta_s^{-1}$ and extend $\bar{(\overline{ab})} = \overline{\bar{a}\bar{b}}$

extends to product $\bar{(ab)} = \bar{a}\bar{b} \Rightarrow \overline{\delta_{s_1 s_2}} = \overline{\delta_{s_1} \delta_{s_2}} = \overline{\delta_{s_1}} \overline{\delta_{s_2}} = \delta_{s_1}^{-1} \delta_{s_2}^{-1}$
 $= (\delta_{s_2} \delta_{s_1})^{-1} = \delta_{s_2 s_1}^{-1} = \delta_{(s_1 s_2)^{-1}} \Rightarrow$ on standard basis $\bar{\delta}_x := (\delta_{x^{-1}})^{-1}$.

Def (KL-basis) $\{b_x \mid x \in W\}$ basis of H satisfying

① (self-duality) $\bar{b}_x = b_x$

② (degree bound) $b_x = \delta_x + \sum_{y < x} h_{y,x} \cdot \delta_y$, for some $h_{y,x} \in v \mathbb{Z}[v]$

Prop KL-basis is unique. Say suppose $\{b_x \mid x \in W\}$ & $\{c_x \mid x \in W\}$ are KL-basis.

\Rightarrow write $b_x = \delta_x + \sum_{y < x} h_{y,x} \delta_y, c_x = \delta_x + \sum_{y < x} g_{y,x} \delta_y$. as in algebra

$\Rightarrow \sum_{y < x} (h_{y,x} - g_{y,x}) \delta_y = b_x - c_x = \overline{b_x - c_x} = \sum_{y < x} (h_{y,x} - g_{y,x}) \cdot \bar{\delta}_y$

~~pick~~ pick ~~highest~~ order y s.t. $h_{y,x} - g_{y,x} \neq 0$. One has $\bar{\delta}_y = \delta_y^{-1} = \delta_y + \sum_{z < y} a_z \delta_z$

So by comparing the highest term on both side

$\Rightarrow h_{y,x} - g_{y,x} = h_{y,x} - g_{y,x} \in v^{-1} \mathbb{Z}[v, v^{-1}] \Rightarrow$ must be zero.

Now one pick $b_s = \delta_s + v$ for $s \in S$. The construction of KL basis explicitly by examples.

Coxeter cplx of (W, S) : Coxeter grp of rank $n = |S| \xrightarrow{\text{option}} \text{in } \mathbb{A}^{n-1}$.

① Take one copy of $\Delta_w^{(n-1)}$ -simplex for each w .

② $\forall w \in W, s \in S$, glue Δ_w to Δ_{ws} via S -gluing
act on right.

G_2 (Type A_2)

$$W = \langle s, t \mid s^2 = t^2 = \text{Id}, sts = tst \rangle$$

quite product rule.

$$\delta_x b_s = \begin{cases} \delta_x s + v \delta_x, & x < x_s \\ \delta_x s + v^{-1} \delta_x, & x > x_s \end{cases}$$

$$\begin{cases} b_{st} = \delta_{st} = 1 \\ b_s = \delta_s + v \\ b_t = \delta_t + v \end{cases}$$



try $b_t \cdot b_s$
 $b_t = \delta_t + v$

$(\rightarrow) b_s$

$$= \delta_{ts} + v \delta_t + v \delta_s + v^2 = b_{ts}$$

by def holds self-dual

try $b_t s \cdot b_t$

$(\rightarrow) b_t$

$$= \delta_{tst} + v(\delta_{st} + \delta_{ts}) + v^2(\delta_s + \delta_t) + v^3 + v$$

$$\Rightarrow b_t s \cdot b_t - b_t = \delta_{tst} + v(\delta_{st} + \delta_{ts}) + v^2(\delta_s + \delta_t) + v^3$$

$b_t b_s b_t - b_s$

\Rightarrow $b_t s t$ by self-dual $b_s b_t b_s - b_s = b_t b_s b_t - b_t$

Same way $\Rightarrow b_{st} = \delta_{st} + v(\delta_t + \delta_s) + v^2$

$$\Rightarrow \begin{pmatrix} 1 \\ b_s \\ b_t \\ b_{st} \\ b_{ts} \\ b_{sts} \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & v & 0 & 0 & 0 \\ 0 & 0 & 1 & v & 0 & 0 \\ & & & & & 1 \end{pmatrix} \begin{pmatrix} 1 \\ \delta_s \\ \delta_t \\ \delta_{st} \\ \delta_{ts} \\ \delta_{sts} \end{pmatrix}$$

What forms the relations in $W = A_2$ is

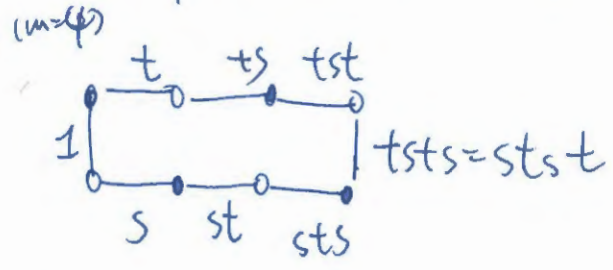
$$\begin{cases} b_s b_s = (v + v^{-1}) b_s \\ b_s b_t b_s - b_s = b_t b_s b_t - b_t \end{cases}$$

in KL basis!

$$\Rightarrow \begin{pmatrix} b_{sts} \\ b_{st} \\ b_{ts} \\ b_s \\ b_t \\ 1 \end{pmatrix} = \begin{bmatrix} 1 & v & v & v^2 & v^2 & v^3 \\ 0 & 1 & 0 & v & v & v^2 \\ 0 & 0 & 1 & v & v & v \\ 0 & 0 & 0 & 1 & 0 & v \\ 0 & 0 & 0 & 0 & 1 & v \\ 0 & 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{pmatrix} \delta_{sts} \\ \delta_{st} \\ \delta_{ts} \\ \delta_s \\ \delta_t \\ 1 \end{pmatrix}$$

Kazhdan-Lusztig ~~form~~ polynomials.

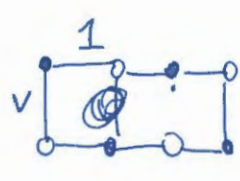
Ex (Type $I_2(m)$) $\begin{matrix} s & \xrightarrow{m} & t \end{matrix}$ means $\delta_{st} = s^2 = t^2 = (st)^m$



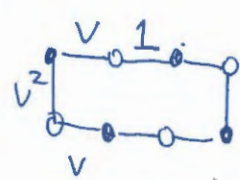
Start with $\begin{cases} b_{st} = \delta_{st} = 1 \\ b_s = \delta_s + v \\ b_t = \delta_t + v \end{cases}$



try $b_t \cdot b_s$
 $b_t = \delta_t + v$



$() \circ b_s$

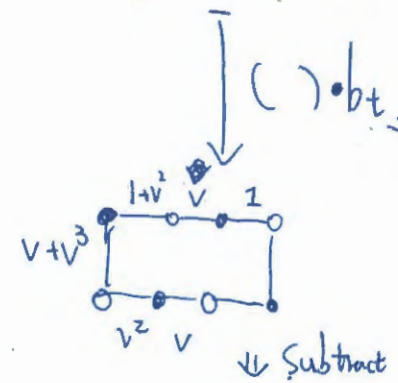


self adj by def $b_{ts} = \delta_{ts} + v(\delta_t + \delta_s) + v^2$

In the symm way

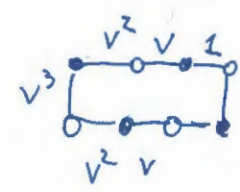
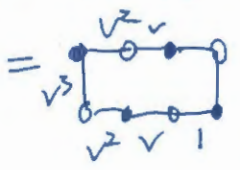


$b_{st} = b_s \cdot b_t = \delta_{st} + v(\delta_s + \delta_t) + v^2$

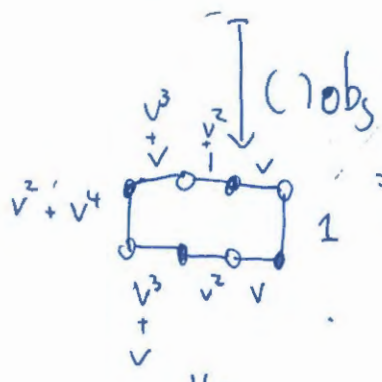


$() \circ b_t = \delta_{st} + v(\delta_{ts} + \delta_{st}) + v^2(\delta_t + \delta_s) + v^3 + [\delta_t + v] b_t$

$b_{ts} = b_{st} \circ b_s - b_s$



$b_{tst} = b_{ts} \cdot b_t - b_t$



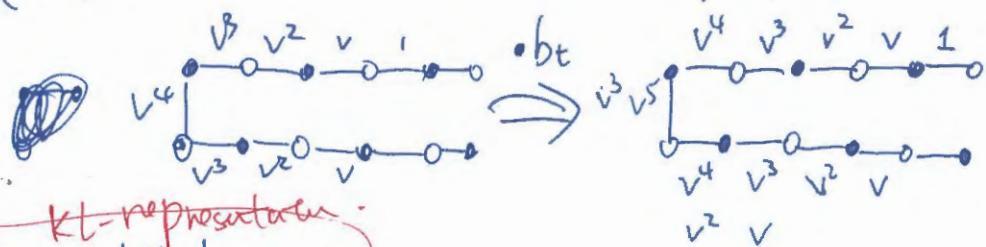
$\delta_{tsts} + v(\delta_{tst} + \delta_{sts}) + v^2(\delta_{ts} + \delta_{st}) + (v^3)(\delta_t + \delta_s) + (v^4) b_{ts}$

1	v	v	v ²	v ²	v ³	v ³	v ⁴	δ_{stst}
1	0	v	v	v ²	v ²	v ³		δ_{tst}
1	v	v	v ²	v ²	v ³			δ_{sts}
		1	0	v	v	v ²		δ_{ts}
			1	0	v	v ²		δ_{st}
				1	0	v		δ_t
					1	v		δ_s
							1	1

In fact the only KL polynomial is "1", up to the difference of orders. On the other hand, [Ref] (Pol 1999)

$b_{tsts} = b_{st} \circ b_s - b_{ts}$

$I_2(m \geq 5)$

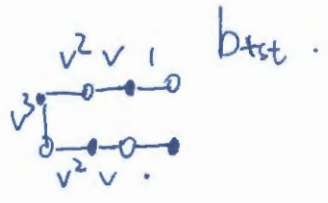


$btsts \cdot bt$

Let's use KL-basis to represent itself:

- $b_{12} = 1$
- b_s
- b_t

subtract



~~btstst~~

$$btstst = btsts \cdot bt - btst$$

- $b_{st} = b_s \cdot b_t$
- $b_{ts} = b_t \cdot b_s$

- $b_{sts} = b_{st} \cdot b_s - b_s = b_s b_t b_s - b_s$

$$b_{tst} = b_{ts} \cdot b_t - b_t = b_t b_s b_t - b_t$$

- $b_{stst} = b_{sts} \cdot b_t - b_{st} = (b_s b_t b_s - b_s) b_t - b_s b_t = b_s b_t b_s b_t - b_s b_t = -2 b_s b_t$

$$b_{tsts} = b_{tst} \cdot b_s - b_{ts} = b_t b_s b_t b_s - 2 b_t b_s$$

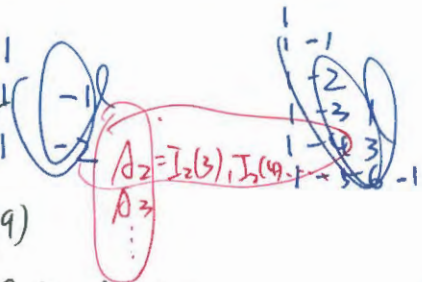
$$b_{ststst} = b_{stst} \cdot b_s - b_{sts} = b_s b_t b_s b_t b_s - 3 b_s b_t b_s + b_s$$

- $b_{ststst} = b_{ststst} \cdot b_t - b_{stst} = b_s b_t b_s b_t b_s b_t - 4 b_s b_t b_s b_t + 3 b_s b_t b_t$

$$b_{tststst} = b_{tstst} \cdot b_s - b_{tsts} = b_t b_s b_t b_s b_t b_s - 4 b_t b_s b_t b_s + 3 b_t b_s b_s$$

- $b_{stststst} = b_{stststst} \cdot b_s - b_{ststst} = b_s b_t b_s b_t b_s b_t b_s - 5 b_s b_t b_s b_t b_s + 6 b_s b_t b_s - b_s$

On the other hand,



Thm (Polol 999)

Any polynomial $P \in \mathbb{N}_{\geq 0}[V]$ w/ const. = 1. appears in KL-polynomials for some S_n .

	1	2	3	4	...
$n-1$	1	-1			
n	1	-2			
$n+1$	1	-3	1		
	1	-4	3		
	1	-5	6	1	
	1	-6	10	-2	
	1	-7	15	-8	1
	1	-8	21	-18	1
	1	-9	28	-33	9
	1	-10	36	-54	27

$$P_{n+1} = P_n - 9P_{n-1}$$

The Standard form on H . Let

$$H \xrightarrow{\delta_x} H \xrightarrow{\delta_x^{-1}} H \quad \overline{ab} = \overline{a} \overline{b}$$

extending

$$H \xrightarrow{\omega} H \quad \omega(ab) = \omega(b)\omega(a)$$

Let $H \xrightarrow{\varepsilon} \mathbb{Z}\langle v, v^\dagger \rangle$
 trace. $\delta_{id} \mapsto 1$
 $\delta_x \mapsto 0, x \neq id$.

Define $H \times H \xrightarrow{(-, -)} \mathbb{Z}\langle v, v^\dagger \rangle \oplus \mathbb{R}$
 $(a, b) \mapsto \varepsilon(\omega(a) \cdot b)$

Prop $(x, y) = (x, \omega(a)y), a \in H$
 $(v^\dagger x, y) = v^\dagger(x, y)$
 $(x, v^\dagger y) = v^\dagger(x, y)$
($v^\dagger x, y = \overline{v}(x, y)$)
($x, v^\dagger y = v(x, y)$)

Prop (orthogonality / bar-orthogonality)
 $\{\delta_x | x \in W\}$ & $\{\delta_x | x \in W\}$ are dual.
 $(\delta_x, \delta_y) = \begin{cases} 1 & x=y \\ 0 & \text{otherwise} \end{cases}$

Cor $\because b\omega(b_s) = b_s$
 $\because (b_s x, y) = (x, b_s y)$
 b_s 's are "self-adj."

Thm (asymptotic orthogonality) For KL-basis,
 $(b_x, b_y) \in \begin{cases} 1 + v\mathbb{Z}\langle v \rangle, & x=y \\ v\mathbb{Z}\langle v \rangle, & \text{otherwise} \end{cases}$

ε & $(-,-)$ are natural for $\mathbb{Z}\langle W \rangle$

KL-basis is the correct choice of basis for v -deformation of $\mathbb{Z}\langle W \rangle$

Deodhar's formula. $W = S_3 = A_2, S = \{s, t\}$. longest element $x = sts$

e	111	110	101	100	011	010	001	000
x^e	sts	st	id	s	ts	t	s	id
defect(e)	0	1	1	0	1	2	2	3

it could repeat.

$\Rightarrow b_s b_t b_s \neq 0 \quad (b_{s_1} \cdots b_{s_m} = \sum_{e \in x} v^{\text{defect}(e)} \delta_x^e)$ $\text{defect}(e) = \#U(e) - \#D(e)$

$$\delta_{sts} + v\delta_{st} + v + \delta_s + v\delta_{ts} + v^2\delta_t + v^2\delta_s + v^3$$